#### Mathematical Induction Part Two

#### Announcements

- Problem Set 1 due Friday, October 4 at the start of class.
- Problem Set 1 checkpoints graded, will be returned at end of lecture.
  - Afterwards, will be available in the filing cabinets in the Gates Open Area near the submissions box.

The **principle of mathematical induction** states that if for some P(n) the following hold:



*Theorem:* For any natural number *n*,  $\sum_{i=0}^{n-1} 2^i = 2^n - 1$ 

*Proof*: By induction. Let P(n) be

$$P(n) \equiv \sum_{i=0}^{n-1} 2^{i} = 2^{n} - 1$$

For our base case, we need to show P(0) is true, meaning that

$$\sum_{i=0}^{-1} 2^i = 2^0 - 1$$

Since  $2^{0} - 1 = 0$  and the left-hand side is the empty sum, P(0) holds.

For the inductive step, assume that for some  $n \in \mathbb{N}$ , that P(n) holds, so  $\sum_{i=1}^{n-1} e^{i} = e^{n}$ 

$$\sum_{i=0}^{n} 2^{i} = 2^{n} - 1$$

We need to show that P(n + 1) holds, meaning that

$$\sum_{i=0}^{n} 2^{i} = 2^{n+1} - 1$$

To see this, note that

$$\sum_{i=0}^{n} 2^{i} = \left(\sum_{i=0}^{n-1} 2^{i}\right) + 2^{n} = 2^{n} - 1 + 2^{n} = 2(2^{n}) - 1 = 2^{n+1} - 1$$

Thus P(n + 1) holds, completing the induction.

## Induction in Practice

- Typically, a proof by induction will not explicitly state P(n).
- Rather, the proof will describe P(n) implicitly and leave it to the reader to fill in the details.
- Provided that there is sufficient detail to determine
  - what P(n) is,
  - that P(0) is true, and that
  - whenever P(n) is true, P(n + 1) is true,

the proof is usually valid.

*Theorem:* For any natural number  $n, \sum_{i=0}^{n-1} 2^i = 2^n - 1$ 

*Proof*: By induction on *n*. For our base case, if n = 0, note that

$$\sum_{i=0}^{-1} 2^i = 0 = 2^0 - 1$$

and the theorem is true for 0.

For the inductive step, assume that for some *n* the theorem is true. Then we have that

$$\sum_{i=0}^{n} 2^{i} = \sum_{i=0}^{n-1} i + 2^{n} = 2^{n} - 1 + 2^{n} = 2(2^{n}) - 1 = 2^{n+1} - 1$$

so the theorem is true for n + 1, completing the induction.

#### Variations on Induction: **Starting Later**

# Induction Starting at 0

- To prove that P(n) is true for all natural numbers greater than or equal to 0:
  - Show that P(0) is true.
  - Show that for any  $n \ge 0$ , that  $P(n) \rightarrow P(n + 1)$ .
  - Conclude P(n) holds for all natural numbers greater than or equal to 0.

# Induction Starting at k

- To prove that P(n) is true for all natural numbers greater than or equal to k:
  - Show that  $P(\mathbf{k})$  is true.
  - Show that for any  $n \ge k$ , that  $P(n) \rightarrow P(n + 1)$ .
  - Conclude P(n) holds for all natural numbers greater than or equal to k.
- Pretty much identical to before, except that the induction begins at a later point.

## Convex Polygons

• A **convex polygon** is a polygon where, for any two points in or on the polygon, the line between those points is contained within the polygon.



### Useful Fact

• **Theorem:** Any line drawn through a convex polygon splits that polygon into two convex polygons.



# Summing Angles

- Interesting fact: the sum of the angles in a convex polygon depends only on the number of vertices in the polygon, not the shape of that polygon.
- **Theorem:** For any convex polygon with *n* vertices, the sum of the angles in that polygon is  $(n 2) \cdot 180^{\circ}$ .
  - Angles in a triangle add up to 180°.
  - Angles in a quadrilateral add up to 360°.
  - Angles in a pentagon add up to 540°.

*Theorem:* The sum of the angles in any convex polygon with *n* vertices is  $(n - 2) \cdot 180^{\circ}$ .

*Proof:* By induction. Let P(n) be "all convex polygons with *n* vertices have angles that sum to  $(n - 2) \cdot 180^{\circ}$ ." We will prove P(n) holds for all  $n \in \mathbb{N}$  where  $n \ge 3$ . As a base case, we prove P(3): the sum of the angles in any convex polygon with three vertices is 180°. Any such polygon is a triangle, so its angles sum to 180°.

For the inductive step, assume for some  $n \ge 3$  that P(n) holds and all convex polygons with n vertices have angles that sum to  $(n-2) \cdot 180^{\circ}$ . We prove P(n+1), that the sum of the angles in any convex polygon with n+1 vertices is  $(n-1) \cdot 180^{\circ}$ . Let A be an arbitrary convex polygon with n+1 vertices. Take any three consecutive vertices in A and draw a line from the first to the third, as shown here:



The sum of the angles in *A* is equal to the sum of the angles in triangle *B* (180°) and the sum of the angles in convex polygon *C* (which, by the IH, is  $(n - 2) \cdot 180^\circ$ ). Therefore, the sum of the angles in *A* is  $(n-1) \cdot 180^\circ$ . Thus P(n + 1) holds, completing the induction.

### A Different Proof Approach



# Using Induction

- Many proofs that work by induction can be written non-inductively by using similar arguments.
- Don't feel that you *have* to use induction; it's one of many tools in your proof toolbox!

#### Variations on Induction: **Bigger Steps**

## Subdividing a Square



For what values of *n* can a square be subdivided into *n* squares?

## The Key Insight

# The Key Insight

- If we can subdivide a square into n squares, we can also subdivide it into n + 3 squares.
- Since we can subdivide a bigger square into 6, 7, and 8 squares, we can subdivide a square into n squares for any  $n \ge 6$ :
  - For multiples of three, start with 6 and keep adding three squares until *n* is reached.
  - For numbers congruent to one modulo three, start with 7 and keep adding three squares until *n* is reached.
  - For numbers congruent to two modulo three, start with 8 and keep adding three squares until *n* is reached.

- Theorem: For any  $n \ge 6$ , it is possible to subdivide a square into n squares.
- *Proof:* By induction. Let P(n) be "a square can be subdivided into n squares." We will prove P(n) holds for all  $n \ge 6$ .
  - As our base cases, we prove P(6), P(7), and P(8), that a square can be subdivided into 6, 7, and 8 squares. This is shown here:

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For the inductive step, assume that for some  $n \ge 6$  that P(n) is true and a square can be subdivided into n squares. We prove P(n + 3), that a square can be subdivided into n + 3 squares. To see this, obtain a subdivision of a square into n squares. Then, choose a square and split it into four equal squares. This removes one of the n squares and adds four more, so there are now n + 3 total squares. Thus P(n + 3) holds, completing the induction.

## **Generalizing Induction**

- When doing a proof by induction:
  - Feel free to use multiple base cases.
  - Feel free to take steps of sizes other than one.
- Just be careful to make sure you cover all the numbers you think that you're covering!

#### Variations on Induction: Complete Induction

#### An Observation





## An Observation

- In a proof by induction, the inductive step works as follows:
  - Assume that for some particular *n* that *P*(*n*) is true.
  - Prove that P(n + 1) is true.
- Notice: When trying to prove P(n + 1), we already know P(0), P(1), P(2), ..., P(n) but only assume P(n) is true.
- Why are we discarding all the intermediary results?

# **Complete Induction**

- If the following are true:
  - P(0) is true, and
  - If P(0), P(1), P(2), ..., P(n) are true, then P(n+1) is true as well.
- Then P(n) is true for all  $n \in \mathbb{N}$ .
- This is called the **principle of complete induction** or the **principle of strong induction**.
  - (A note: this also works starting from a number other than 0; just modify what you're assuming appropriately.)

### Proof by Complete Induction

- State that your proof works by complete induction.
- State your choice of P(n).
- Prove the base case: state what P(0) is, then prove it using any technique you'd like.
- Prove the inductive step:
  - State that for some arbitrary  $n \in \mathbb{N}$  that you're assuming P(0), P(1), ..., P(n) (that is, P(n') for all natural numbers  $0 \le n' \le n$ .)
  - State that you are trying to prove P(n + 1) and what P(n + 1) means.
  - Prove P(n + 1) using any technique you'd like.

### Example: Polygon Triangulation



# Polygon Triangulation

- Given a convex polygon, an elementary triangulation of that polygon is a way of connecting the vertices with lines such that
  - No two lines intersect, and
  - The polygon is converted into a set of triangles.
- Question: How many lines do you have to draw to elementarily triangulate a convex polygon?

![](_page_29_Figure_1.jpeg)

![](_page_30_Picture_1.jpeg)

![](_page_31_Figure_1.jpeg)

### Some Observations

- Every elementary triangulation of the same convex polygon seems to require the same number of lines.
- The number of lines depends on the number of vertices:
  - 5 vertices: 2 lines
  - 6 vertices: 3 lines
  - 8 vertices: 5 lines
- Conjecture: Every elementary triangulation of an *n*-vertex convex polygon requires *n* – 3 lines.

![](_page_33_Figure_1.jpeg)

Theorem: Every elementary triangulation of a convex polygon with n vertices requires n - 3 lines.

*Proof:* By complete induction. Let P(n) be "every elementary triangulation of a convex polygon requires n-3 lines." We prove P(n) holds for all  $n \ge 3$ . As a base case, we prove P(3): elementarily triangulating a convex polygon with three vertices requires no lines. Any polygons with three vertices is a triangle, so any elementary triangulation of it will have no lines.

For the inductive step, assume for some  $n \ge 3$  that P(3), P(4), ..., P(n) are true. This means any elementary triangulation of an n'-vertex convex polygon, where  $3 \le n' \le n$ , uses n'-3 lines. We prove P(n+1): any elementary triangulation of any (n+1)-vertex convex polygon uses n-2 lines.

Let *A* be an arbitrary convex polygon with n+1 vertices. Pick any elementary triangulation of *A* and select an arbitrary line in that triangulation. This line splits *A* into two smaller convex polygons *B* and *C*, which are also triangulated. Let *k* be the number of vertices in *B*, meaning *C* has (n+1)-k+2 = n-k+3 vertices. By our inductive hypothesis, any triangulations of *B* and *C* must use k-3 and n-k lines, respectively. Therefore, the total number of lines in the triangulation of *A* is n-k+k-3+1 = n-2. Thus P(n+1) holds, completing the induction.

# Using Complete Induction

- When is it appropriate to use complete induction in contrast to standard induction?
- Depends on the proof approach:
  - Typically, standard induction is used when a problem of size n + 1 is reduced to a simpler problem of size n.
  - Typically, complete induction is used when the problem of size n + 1 is split into multiple subproblems of unknown but smaller sizes.
- It is never "wrong" to use complete induction. It just might be unnecessary. We suggest writing drafts of your proofs just in case.

## Summary

- Induction doesn't have to start at 0. It's perfectly fine to start induction later on.
- Induction doesn't have to take steps of size 1. It's not uncommon to see other step sizes.
- Induction doesn't have to have a single base case.
- Complete induction lets you assume all prior results, not just the last result.

### Next Time

#### Graphs

- Representing relationships between objects.
- Connectivity in graphs.
- Planar graphs.